

1. Chapter I: #10ab.

**ANS:** This is a test to see if you can parse a definition precisely. Recall that a function  $f : A \rightarrow B$  is a subset  $f \subset A \times B$  such that for every  $a \in A$ , there is a unique  $b \in B$  such that  $(a, b) \in f$ . If either  $A$  or  $B$  is empty, then so is  $A \times B$ . Hence there is only one subset of  $A \times B$  — namely the empty set  $\emptyset$ . The only question is whether or not the empty set is a function.

- (a) If  $A$  is nonempty and  $B = \emptyset$ , then given  $a \in A$ , there can be no  $(a, b) \in A \times B = \emptyset$ , so there are no functions  $f : A \rightarrow \emptyset$ .
- (b) On the other hand, if  $A = \emptyset$ , then whether or not  $B$  is empty, the condition for all  $a \in A$  there is a unique  $b \in B$  such that  $(a, b) \in \emptyset$  is vacuously satisfied. Hence the empty set is a function, and the only function, from  $\emptyset$  to  $B$ .

2. Chapter II: #11.

**ANS:** Here we have to show that if  $a > 1$ , then  $\{a, a^2, a^3, \dots\}$  is not bounded. Ok, suppose not. Then there is  $x \in \mathbf{R}$  such that  $a^n \leq x$  for all  $n \in \mathbf{N}$ .

Next we turn to the hint. We'll show that

$$\left(1 + \frac{1}{n}\right)^n \geq 2 \quad \text{for all } n \in \mathbf{N} \tag{†}$$

using induction.<sup>1</sup> Let  $A$  be the subset of  $\mathbf{N}$  for which (†) holds. Clearly  $1 \in A$  and if  $n \in A$ , then

$$\begin{aligned} \left(1 + \frac{1}{n}\right)^{n+1} &= \left(1 + \frac{1}{n}\right)^n \left(1 + \frac{1}{n}\right) \\ &\geq 2 \cdot 1 = 2. \end{aligned}$$

(Here we've used  $n \in A$  and  $(1 + \frac{1}{n}) \geq 1$ .) This shows (†) holds for all  $n$ .

Next I claim that

$$2^k \geq k \quad \text{for all } k \in \mathbf{N}. \tag{‡}$$

Again, we'll use induction. Let  $A$  be the set of  $k$  for which (‡) holds. Clearly  $1 \in A$ . Suppose  $n \in A$ . Then

$$2^{n+1} \geq 2^n(2) \geq 2n = n + n \geq n + 1.$$

Thus (‡) holds for all  $k$ . But there is a  $k > x$  (by LUB 1). Since  $a - 1 > 0$ , there is a  $n \in \mathbf{N}$  such that  $\frac{1}{n} < a - 1$  and

$$a < \left(1 + \frac{1}{n}\right).$$

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<sup>1</sup>Alternately, you could use the result we proved in lecture that  $(1 + x)^n \leq 1 + nx$  provided  $x \geq -1$ . You can't for example, use the binomial theorem as we haven't proved it. Of course, you could prove it and then use it.

But then

$$\begin{aligned} a^{kn} &> \left( \left( 1 + \frac{1}{n} \right)^n \right)^k \\ &\geq 2^k \\ &\geq k \\ &> x. \end{aligned}$$

But this contradicts our choice of  $x$ . This finishes the proof.

3. Chapter II: #13.

**ANS:** Since each  $S_i$  is nonempty and bounded above, each set has a least upper bound  $s_i$ . Define

$$S_1 + S_2 = \{ x + y : x \in S_1 \text{ and } y \in S_2 \}.$$

We are supposed to show that  $\text{lub}(S_1 + S_2) = s_1 + s_2$ . But if  $x \in S_1$  and  $y \in S_2$ , then

$$x + y \leq s_1 + s_2.$$

Hence  $S_1 + S_2$  is bounded above (as well as nonempty). Hence  $S_1 + S_2$  at least has an least upper bound. Since  $s_1 + s_2$  is an upper bound, it will suffice to see that  $s_1 + s_2 - \epsilon$  is not an upper bound for any  $\epsilon > 0$ . But  $s_1 - \epsilon/2$  can't be an upper bound for  $S_1$ . Thus there is a  $t_1 \in S_1$  such that  $t_1 > s_1 - \epsilon/2$ . Similarly, there is a  $t_2 \in S_2$  such that  $t_2 > s_2 - \epsilon/2$ . But now we have  $t_1 + t_2 \in S_1 + S_2$  and

$$t_1 + t_2 > s_1 + s_2 - \epsilon.$$

Thus  $s_1 + s_2 - \epsilon$  is not an upper bound and we're done.