

# Math 56 Compu & Expt Math, Spring 2013: Homework 4

due 10am Thursday April 25th

*Exploring beautiful properties and applications of things Fourier.*

1. Fourier series theory. As in lecture, let  $\|\cdot\|$  be the  $L_2$ -norm on  $(0, 2\pi)$ .
  - (a) What is  $\|e^{inx}\|$  for any  $n \in \mathbb{Z}$ ?
  - (b) Find a function orthogonal to  $f(x) = x$  on  $0 \leq x < 2\pi$ .
  - (c) Combine the Fourier series for the  $2\pi$ -periodic function defined by  $f(x) = x$  for  $0 \leq x < 2\pi$  that we computed on the worksheet with Parseval's relation to evaluate  $\sum_{m=1}^{\infty} m^{-2}$
  - (d) Compute the Fourier series for the  $2\pi$ -periodic function defined by  $f(x) = x^2$  in  $-\pi \leq x < \pi$ . [Hint: shift the domain of integration to a convenient one.] Comment on how  $\hat{f}_m$  decays for this continuous (but not  $C^1$ ) function compared to the discontinuous function in (c).
  - (e) Take the expression for a general  $f$  written as a Fourier series. By taking the derivative of both sides (you may assume that you can pass the derivative through the sum), prove that if  $|f'|$  is bounded,  $|\hat{f}_m| = O(1/|m|)$ .
  - (f) Use the previous idea to prove a bound on the decay of Fourier coefficients when  $f$  has  $k$  bounded derivatives. What bound follows if  $f \in C^\infty$  (arbitrarily smooth)? Can you give this a name?
2. Getting to know your DFT. Use numerical exploration followed by proof (each proof is very quick):
  - (a) Produce a color image of the real part of the DFT matrix  $F$  for  $N = 256$ . Explain one of the *curves* seen.
  - (b) What is  $F^2$ ? [careful: matrix product, also don't forget the 0-indexing]. What does  $F^2$  do to a vector? (this should be very simple!) Now, for general  $N$ , prove your claim [Hint: use  $\omega$ ]
  - (c) What then is  $F^4$ ? Prove this.
  - (d) What are the eigenvalues of  $F$ ? Use your previous result to prove this.
  - (e) What is the condition number of  $F$ ? Prove this using a result from lecture.
3. The power of trigonometric interpolation, i.e. using just a few samples of a periodic function to reconstruct the function *everywhere*. (Applications to modeling data, etc.)
  - (a) Let's interpolate  $f(x) = e^{\sin x}$ . For  $N = 40$ , by using the  $1/N$ -weighted samples at the nodes  $x_j = 2\pi j/N$ ,  $j = 0, \dots, N-1$ , and `fft`, find  $\tilde{f}_m$ . Plot their magnitudes on a log vertical scale vs  $m = 0, \dots, N-1$ . Relate to  $\#1(f)$ . By what  $|m|$  have the coefficients decayed to  $\varepsilon_{\text{mach}}$  times the largest? (This is the effective band-limit of the function at this tolerance.)
  - (b) Using  $\tilde{f}_m$  as good approximations to the true Fourier coefficients in  $-N/2 < m < N/2$ , plot the "interpolant" given by this truncated Fourier series, on the fine grid `0:1e-3:2*pi`. Overlay the samples  $Nf_j = f(x_j)$  as blobs. [Hint: debug until the interpolant passes through the samples]
  - (c) By looping over the above for different  $N$ , make a labeled semi-log plot of the maximum error (taken over the fine grid) between the interpolant and  $f$ , vs  $N$ , for even  $N$  between 2 and 40. At what  $N$  is convergence to  $\varepsilon_{\text{mach}}$  reached? (Pretty amazing, eh?) Relate to (a).
4. Let's prove that, amongst all trigonometric polynomials of degree at most  $N/2$ , the  $N/2$ -truncated Fourier series for  $f$  is the *best approximation* to  $f$  in the  $L_2(0, 2\pi)$  norm.

- (a) Any trig. poly. can be written  $\sum_{|n| \leq N/2} (\hat{f}_n + c_n) e^{inx}$  for some coefficient “deviations”  $c_n$ . Write the squared  $L_2$ -norm of the function which is the difference of the above and the true  $f$ .
- (b) Expand out to four terms, use the definition of  $\hat{f}_n$ , then expand further, and cancel stuff to leave all the  $c_n$  dependence in a sum of squares. Your  $\text{\LaTeX}$  only needs to show 3-4 key steps. The proof is now easy.
- (c) From that, extract an expression for the square of this (best) error.

5. How fast do things run?

- (a) Consider the matvec problem computing  $\mathbf{y} = \mathbf{A}\mathbf{x}$  for  $\mathbf{A}$  an  $n$ -by- $n$  matrix. Write a little code using random  $\mathbf{A}$  and  $\mathbf{x}$  for  $n = 4000$ , which measures the runtime of Matlab’s native  $\mathbf{A}*\mathbf{x}$  and your own naive double loop to compute the same thing. Express your answers in “flops” (flop per sec), and give the ratio. Marvel at how fast the built-in library is.
- (b) Make a plot showing how many microseconds ( $\mu\text{s} = 10^{-6}$  s) it takes to do a FFT of a random vector of length  $n$ , varying  $n$  over integers between 8100 to 8200 (use at least 100 repetitions for each  $n$  to get runtimes that are long enough to measure accurately). Overlay on your plot the reciprocal of the number of prime factors of  $n$ , scaled vertically so as to have the same max value. Which  $n$  has the fastest FFT (why?). How many times slower is the slowest?